

Proof Appendix of “The Substitution Vanishes”

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Abstract. This is the proof appendix of “The Substitution Vanishes”. It contains the proofs of Lemma 13 (see Lemma 3), Theorem 14 (see Theorem 4), Lemma 20 (see Lemma 8), and Theorem 21 (see Theorem 9).

Appendix

Lemma 1. Let $M = (F, \{sub^{(n+1)}\}, C, \Pi, R_1, R_2, r_{in})$ be an n -sntt.

$$\underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, sub\ e_0 \dots e_n)) = \underline{tree}(e_0)[\Pi_i / \underline{tree}(e_i)]$$

for every $e_0, \dots, e_n \in E_{C,n}$ such that if there is a subtree of the shape $e'_0[z_i \rightsquigarrow e'_i]$ for some $e'_0 \in E_{C,n}(Z_n)$ and $e'_1, \dots, e'_n \in E_{C,n}$ in e_0 , then e'_0 does not contain any constructor of Π .

Proof. We just write $nf(r)$ when we mean $nf(\Rightarrow_{R_1 \cup R_2}, r)$.

(i) First let $e_0 = \Pi_j$ for some $j \in [n]$.

$$\begin{aligned} & \underline{tree}(nf(sub\ \Pi_j\ e_1 \dots e_n)) \\ = & \text{(by reduction and because } e_j \text{ is in normal form)} \\ & \underline{tree}(e_j) \\ = & \text{(by definition of } \underline{tree} \text{ and substitution)} \\ & \underline{tree}(\Pi_j)[\Pi_i / \underline{tree}(e_i)] \end{aligned}$$

(ii) Let $e_0 = c\ e'_1 \dots e'_m$ for some $m \in \mathbb{N}$, $c \in (C - \Pi)^{(m)}$, and $e'_1, \dots, e'_m \in E_{C,n}$.

$$\begin{aligned} & \underline{tree}(nf(sub\ (c\ e'_1 \dots e'_m)\ e_1 \dots e_n)) \\ = & \text{(by reduction)} \\ & \underline{tree}(nf((c\ (sub\ e'_1\ z_1 \dots z_n) \dots (sub\ e'_m\ z_1 \dots z_n))[z_i \rightsquigarrow e_i])) \\ = & \text{(by normal form and definition of } \underline{tree} \text{ because } e_i \text{ is in normal form)} \\ & (c\ \underline{tree}(nf(sub\ e'_1\ z_1 \dots z_n)) \dots \\ & \quad \underline{tree}(nf(sub\ e'_m\ z_1 \dots z_n)))[z_i / \underline{tree}(e_i)] \\ = & \text{(by induction hypothesis [} m \text{ times])} \end{aligned}$$

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$$\begin{aligned}
& (c \text{ tree}(e'_1)[\Pi_i/\text{tree}(z_i)] \dots \text{tree}(e'_m)[\Pi_i/\text{tree}(z_i)])[z_i/\text{tree}(e_i)] \\
= & \text{ (by substitution and definition of } \underline{\text{tree}}) \\
& \underline{\text{tree}}(c e'_1 \dots e'_m)[\Pi_i/z_i][z_i/\text{tree}(e_i)] \\
= & \text{ (by substitution because } \underline{\text{tree}}(c e'_1 \dots e'_m) \in T_C) \\
& \underline{\text{tree}}(c e'_1 \dots e'_m)[\Pi_i/\text{tree}(e_i)]
\end{aligned}$$

(iii) Finally, let $e_0 = e'_0[z_i \rightsquigarrow e'_i]$ for some $e'_0 \in E_{C,n}(Z_n)$ and $e'_1, \dots, e'_n \in E_{C,n}$.

$$\begin{aligned}
& \underline{\text{tree}}(\text{nf}(\text{sub } e'_0[z_i \rightsquigarrow e'_i] e_1 \dots e_n)) \\
= & \text{ (by reduction)} \\
& \underline{\text{tree}}(\text{nf}(e'_0[z_i \rightsquigarrow \text{sub } e'_i z_1 \dots z_n][z_i \rightsquigarrow e_i])) \\
= & \text{ (by normal form because } e_i \text{ and } e'_0 \text{ are in normal form)} \\
& \underline{\text{tree}}(e'_0[z_i \rightsquigarrow \text{nf}(\text{sub } e'_i z_1 \dots z_n)][z_i \rightsquigarrow e_i]) \\
= & \text{ (by definition of } \underline{\text{tree}}) \\
& \underline{\text{tree}}(e'_0[z_i/\text{tree}(\text{nf}(\text{sub } e'_i z_1 \dots z_n))][z_i/\text{tree}(e_i)]) \\
= & \text{ (by induction hypothesis [} n \text{ times])} \\
& \underline{\text{tree}}(e'_0[z_i/\text{tree}(e'_i)[\Pi_i/\text{tree}(z_i)]][z_i/\text{tree}(e_i)]) \\
= & \text{ (by associativity of substitution} \\
& \quad \text{because } \underline{\text{tree}}(e'_0) \text{ does not contain constructors of } \Pi) \\
& \underline{\text{tree}}(e'_0[z_i/\text{tree}(e'_i)][\Pi_i/\text{tree}(z_i)][z_i/\text{tree}(e_i)]) \\
= & \text{ (by definition of } \underline{\text{tree}}) \\
& \underline{\text{tree}}(e'_0[z_i \rightsquigarrow e'_i][\Pi_i/z_i][z_i/\text{tree}(e_i)]) \\
= & \text{ (by substitution because } \underline{\text{tree}}(e'_0[z_i \rightsquigarrow e'_i]) \in T_C) \\
& \underline{\text{tree}}(e'_0[z_i \rightsquigarrow e'_i][\Pi_i/\text{tree}(e_i)])
\end{aligned}$$

This completes the proof of the lemma. \square

Lemma 2. Let $M = (F, \text{Sub}, C, \Pi, R_1, R_2, r_{in})$ with $\text{Sub} = \{\text{sub}^{(n+1)}\}$ be an n -sntt.

$$\begin{aligned}
& \underline{\text{tree}}(\text{nf}(\Rightarrow_{R_1 \cup R_2}, (\text{sub } r_0 \dots r_n)[x_i/t_i])) \\
= & \underline{\text{tree}}(\text{nf}(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i]))[\Pi_j/\text{tree}(\text{nf}(\Rightarrow_{R_1 \cup R_2}, r_j[x_i/t_i]))]
\end{aligned}$$

for every $m \in \mathbb{N}$, $t_1, \dots, t_m \in T_C$, and $r_0, \dots, r_n \in \text{RHS}(F, \text{Sub}, C, X_m)$.

Proof. We just write $\text{nf}(r)$ when we mean $\text{nf}(\Rightarrow_{R_1 \cup R_2}, r)$.

$$\begin{aligned}
& \underline{\text{tree}}(\text{nf}((\text{sub } r_0 \dots r_n)[x_i/t_i])) \\
= & \text{ (by substitution)} \\
& \underline{\text{tree}}(\text{nf}(\text{sub } r_0[x_i/t_i] \dots r_n[x_i/t_i])) \\
= & \text{ (by normal form)} \\
& \underline{\text{tree}}(\text{nf}(\text{sub } \text{nf}(r_0[x_i/t_i]) \dots \text{nf}(r_n[x_i/t_i])))
\end{aligned}$$

Let $e_i = nf(r_i[x_i/t_i])$ for every $0 \leq i \leq n$. Clearly, $e_i \in E_{C,n}$. Note that whenever there is a subtree of the shape $e'_0[z_i \rightsquigarrow e'_i]$ for some $e'_0 \in E_{C,n}(Z_n)$ and $e'_1, \dots, e'_n \in E_{C,n}$ in e_0 , then e'_0 does not contain any constructor of Π . This property can easily be observed by considering the rules for *sub*. It remains to prove

$$\underline{tree}(nf(\text{sub } e_0 \dots e_n)) = \underline{tree}(e_0)[\Pi_i/\underline{tree}(e_i)] ,$$

which holds by Lemma 1. \square

Lemma 3 (see Lemma 13). Let $M = (F, \text{Sub}, C, \Pi, R_1, R_2, r_{in})$ with $\text{Sub} = \{\text{sub}^{(n+1)}\}$ be an n -sntt, and let $\text{acc}(M) = (\text{acc}(F), C, \text{acc}(R_1), \text{acc}(r_{in}))$ be the n -satt constructed from M by accumulation.

(a) For every $f \in F$ and $t \in T_C$:

$$\underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \text{sub } (f \ t) \ z_1 \dots z_n)) = \underline{tree}(nf(\Rightarrow_{\text{acc}(R_1)}, f \ t \ z_1 \dots z_n)) .$$

(b) For every $k \in \mathbb{N}$, $r \in \text{RHS}(F, \text{Sub}, C, X_k)$, and $t_1, \dots, t_k \in T_C$:

$$\begin{aligned} & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \text{sub } r[x_i/t_i] \ z_1 \dots z_n)) \\ &= \underline{tree}(nf(\Rightarrow_{\text{acc}(R_1)}, nf(\Rightarrow_{R_2}, \underline{\text{sub}} \ r \ z_1 \dots z_n)[x_i/t_i])) . \end{aligned}$$

Proof. We prove the two statements by simultaneous induction. Statement (a) is proved by structural induction on t , so let $t = c \ t_1 \dots t_k$ for some $k \in \mathbb{N}$, $c \in C^{(k)}$, and $t_1, \dots, t_k \in T_C$.

$$\begin{aligned} & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \text{sub } (f \ (c \ t_1 \dots t_k)) \ z_1 \dots z_n)) \\ &= \text{(by reduction)} \\ & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \text{sub } \text{rhs}_{R_1, f, c}[x_i/t_i] \ z_1 \dots z_n)) \\ &= \text{(by Statement (b))} \\ & \underline{tree}(nf(\Rightarrow_{\text{acc}(R_1)}, nf(\Rightarrow_{R_2}, \underline{\text{sub}} \ \text{rhs}_{R_1, f, c} \ z_1 \dots z_n)[x_i/t_i])) \\ &= \text{(by substitution)} \\ & \underline{tree}(nf(\Rightarrow_{\text{acc}(R_1)}, nf(\Rightarrow_{R_2}, \underline{\text{sub}} \ \text{rhs}_{R_1, f, c} \ y_1 \dots y_n)[x_i/t_i] [y_i/z_i])) \\ &= \text{(by definition of } \text{acc}(R_1) \text{ and reduction)} \\ & \underline{tree}(nf(\Rightarrow_{\text{acc}(R_1)}, f \ (c \ t_1 \dots t_k) \ z_1 \dots z_n)) \end{aligned}$$

Statement (b) is proved by induction on r .

(i) First let $r = \Pi_j$ for some $j \in [n]$.

$$\begin{aligned} & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \text{sub } \Pi_j[x_i/t_i] \ z_1 \dots z_n)) \\ &= \text{(by substitution and reduction)} \\ & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, z_j)) \\ &= \text{(by } nf(\Rightarrow_{R_1 \cup R_2}, \cdot)) \end{aligned}$$

$$\begin{aligned}
& \underline{tree}(z_j) \\
= & \text{(by } nf(\Rightarrow_{acc(R_1)}, \cdot) \text{ and substitution)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, z_j[x_i/t_i])) \\
= & \text{(by reduction)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} \Pi_j z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

(ii) Now let $r = f x_j$ for some $f \in F$ and $j \in [k]$.

$$\begin{aligned}
& \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} (f x_j)[x_i/t_i] z_1 \dots z_n)) \\
= & \text{(by substitution)} \\
& \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} (f t_j) z_1 \dots z_n)) \\
= & \text{(by Statement (a))} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, f t_j z_1 \dots z_n)) \\
= & \text{(by substitution)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, (f x_j z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by reduction)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} (f x_j) z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

(iii) Let $r = c r_1 \dots r_m$ for some $m \in \mathbb{N}$, $c \in (C - \Pi)^{(m)}$, and $r_1, \dots, r_m \in RHS(F, Sub, C, X_k)$.

$$\begin{aligned}
& \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} (c r_1 \dots r_m)[x_i/t_i] z_1 \dots z_n)) \\
= & \text{(by substitution)} \\
& \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} (c r_1[x_i/t_i] \dots r_m[x_i/t_i]) z_1 \dots z_n)) \\
= & \text{(by reduction)} \\
& \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, (c (\underline{sub} r_1[x_i/t_i] z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_m[x_i/t_i] z_1 \dots z_n))[z_i \rightsquigarrow z_i])) \\
= & \text{(by } nf(\Rightarrow_{R_1 \cup R_2}, \cdot), \text{ definition of } \underline{tree}, \text{ and substitution)} \\
& c \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_1[x_i/t_i] z_1 \dots z_n)) \dots \\
& \quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_m[x_i/t_i] z_1 \dots z_n)) \\
= & \text{(by induction hypothesis [} m \text{ times])} \\
& c \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_1 z_1 \dots z_n)[x_i/t_i])) \dots \\
& \quad \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_m z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by } nf(\Rightarrow_{acc(R_1)}, \cdot), nf(\Rightarrow_{\overline{R_2}}, \cdot), \text{ definition of } \underline{tree}, \text{ and substitution)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \\
& \quad c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \\
= & \text{(by substitution)} \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \\
& \quad c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n)) [x_i/t_i])) [z_i/z_i]
\end{aligned}$$

$$\begin{aligned}
&= \text{(by definition of } \underline{tree} \text{ and } nf(\Rightarrow_{acc(R_1)}, \cdot)\text{)} \\
&\quad \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \\
&\quad\quad c(\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n)) [x_i/t_i] [z_i \rightsquigarrow z_i])) \\
&= \text{(by substitution and } nf(\Rightarrow_{\overline{R_2}}, \cdot)\text{)} \\
&\quad \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \\
&\quad\quad (c(\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[z_i \rightsquigarrow z_i]) [x_i/t_i])) \\
&= \text{(by reduction)} \\
&\quad \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub}(c r_1 \dots r_m) z_1 \dots z_n)[x_i/t_i])) \\
\text{(iv) Finally, let } r = \underline{sub} r_0 \dots r_n \text{ for some } r_0, \dots, r_n \in RHS(F, Sub, C, X_k)\text{.} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub}(\underline{sub} r_0 \dots r_n)[x_i/t_i] z_1 \dots z_n)) \\
&= \text{(by substitution)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub}(\underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i]) z_1 \dots z_n)) \\
&= \text{(by Lemma 2)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i])) [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, z_i))] \\
&= \text{(by Lemma 2)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) \\
&\quad\quad [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_i[x_i/t_i]))] [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, z_i))] \\
&= \text{(by associativity of substitution)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) \\
&\quad\quad [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_i[x_i/t_i]))] [II_j / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, z_j))] \\
&= \text{(by Lemma 2 [} n \text{ times])} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) \\
&\quad\quad [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_i[x_i/t_i] z_1 \dots z_n))] \\
&= \text{(by induction hypothesis [} n \text{ times])} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) \\
&\quad\quad [II_i / \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i]))] \\
&= \text{(by substitution because } \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) \in TC\text{)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) [II_i / z_i] \\
&\quad\quad [z_i / \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i]))] \\
&= \text{(by } nf(\Rightarrow_{R_1 \cup R_2}, \cdot)\text{ and definition of } \underline{tree}\text{)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_0[x_i/t_i])) [II_i / \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, z_i))] \\
&\quad\quad [z_i / \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i]))] \\
&= \text{(by Lemma 2)} \\
&\quad \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_0[x_i/t_i] z_1 \dots z_n)) \\
&\quad\quad [z_i / \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i]))] \\
&= \text{(by induction hypothesis)}
\end{aligned}$$

$$\begin{aligned}
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \\
& \quad [z_i/\underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i]))] \\
= & \quad (\text{since } \underline{tree} \text{ distributes over substitutions}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \\
& \quad [z_i/nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i])] \\
= & \quad (\text{since } nf(\Rightarrow_{acc(R_1)}, \cdot) \text{ distributes over substitutions}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \\
& \quad [z_i/nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_i z_1 \dots z_n)[x_i/t_i])] \\
= & \quad (\text{by } nf(\Rightarrow_{\overline{R_2}}, \cdot) \text{ and substitution}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[x_i/t_i])) \\
= & \quad (\text{by substitution}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[x_i/t_i]))[z_i/z_i] \\
= & \quad (\text{by definition of } \underline{tree}, nf(\Rightarrow_{acc(R_1)}, \cdot), \text{ and substitution}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[z_i \rightsquigarrow z_i][x_i/t_i])) \\
= & \quad (\text{by } nf(\Rightarrow_{\overline{R_2}}, \cdot)) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, (\underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[z_i \rightsquigarrow z_i])[x_i/t_i])) \\
= & \quad (\text{by reduction}) \\
& \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} (sub r_0 \dots r_n) z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

This completes the proof. \square

Theorem 4 (see Theorem 14). Let M be an n -sntt. Then, $\tau_M = \tau_{acc(M)}$.

Proof. Let $M = (F, \{sub^{(n+1)}\}, C, \Pi, R_1, R_2, r_{in})$ be an n -sntt and $acc(M) = (acc(F), C, acc(R_1), acc(r_{in}))$. Then, for every $t \in T_C$:

$$\begin{aligned}
& \tau_M(t) \\
= & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, r_{in}[x_1/t])) && (\text{Def. } \tau_M) \\
= & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} nf(\Rightarrow_{R_1 \cup R_2}, r_{in}[x_1/t]) \Pi_1 \dots \Pi_n)) && (*) \\
= & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} nf(\Rightarrow_{R_1 \cup R_2}, r_{in}[x_1/t]) z_1 \dots z_n)[z_i/\Pi_i]) \\
= & \underline{tree}(nf(\Rightarrow_{R_1 \cup R_2}, \underline{sub} r_{in}[x_1/t] z_1 \dots z_n)[z_i/\Pi_i]) && (\text{confluence}) \\
= & \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_{in} z_1 \dots z_n)[x_1/t]))[z_i/\Pi_i] && (\text{Lm. 3(b)}) \\
= & \underline{tree}(nf(\Rightarrow_{acc(R_1)}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_{in} \Pi_1 \dots \Pi_n)[x_1/t])) \\
= & \underline{tree}(nf(\Rightarrow_{acc(R_1)}, acc(r_{in}))[x_1/t])) && (\text{Def. } acc(r_{in})) \\
= & \tau_{acc(M)}(t) && (\text{Def. } \tau_{acc(M)})
\end{aligned}$$

The statement (*) holds, because every occurrence of a Π_j is substituted by Π_j . (*) can be formally proved by induction on expressions $e_0 \in E_{C,n}$ such that

if there is a subtree of the shape $e'_0[z_i \rightsquigarrow e'_i]$ for some $e'_0 \in E_{C,n}(Z_n)$ and $e'_1, \dots, e'_n \in E_{C,n}$ in e_0 , then e'_0 does not contain any constructor of Π . Note that $(*)$ does not hold without the outer application of tree. \square

Lemma 5. Let C be a ranked alphabet, and let $e_1, e_2 \in E_{C,n}(Z_n)$ and $z \in Z_n$. If tree(e_1) = tree(e_2), then rel(z, e_1) = rel(z, e_2).

Proof. The proof is trivial and hence omitted. \square

Lemma 6. Let $M = (F, \{sub^{(n+1)}\}, C, \Pi, R_1, R_2, r_{in})$ be an n -sntt, and let $e_0, \dots, e_n \in E_{C^\circ, n}$.

$$\begin{aligned} & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub\ e_0 \dots e_n)) \\ &= \underline{step}(e_0) + \sum_{i=1}^n \underline{rel}(z_i, e_0[\Pi_l/z_l]) \cdot \underline{step}(e_i) \end{aligned}$$

Proof. For brevity, we abbreviate $nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r)$ by just $nf(r)$. We prove the statement by induction on e_0 .

(i) First let $e_0 = \Pi_j$ for some $j \in [n]$.

$$\begin{aligned} & \underline{step}(nf(sub\ \Pi_j\ e_1 \dots e_n)) \\ &= \text{(by reduction and normal form because } e_j \text{ is in normal form)} \\ & \quad \underline{step}(e_j) \\ &= \text{(by definition of } \underline{step} \text{ and } \underline{rel} \text{ and substitution)} \\ & \quad \underline{step}(\Pi_j) + \sum_{i=1}^n \underline{rel}(z_i, \Pi_j[\Pi_l/z_l]) \cdot \underline{step}(e_i) \end{aligned}$$

(ii) Now let $e_0 = c\ e'_1 \dots e'_m$ for some $m \in \mathbb{N}$, $c \in (C - \Pi)^{(m)}$, and $e'_1, \dots, e'_m \in E_{C^\circ, n}$. Let $a = 1$ if $c = \circ$ and $a = 0$ otherwise.

$$\begin{aligned} & \underline{step}(nf(sub\ (c\ e'_1 \dots e'_m)\ e_1 \dots e_n)) \\ &= \text{(by reduction)} \\ & \quad \underline{step}(nf((c\ (sub\ e'_1\ z_1 \dots z_n) \dots (sub\ e'_m\ z_1 \dots z_n))[z_i \rightsquigarrow e_i])) \\ &= \text{(by normal form because } e_i \text{ is in normal form)} \\ & \quad \underline{step}(nf(c\ (sub\ e'_1\ z_1 \dots z_n) \dots (sub\ e'_m\ z_1 \dots z_n))[z_i \rightsquigarrow e_i]) \\ &= \text{(by definition of } \underline{step}) \\ & \quad \underline{step}(nf(c\ (sub\ e'_1\ z_1 \dots z_n) \dots (sub\ e'_m\ z_1 \dots z_n))) + \\ & \quad + \sum_{i=1}^n \underline{rel}(z_i, nf(c\ (sub\ e'_1\ z_1 \dots z_n) \dots (sub\ e'_m\ z_1 \dots z_n))) \cdot \underline{step}(e_i) \\ &= \text{(by normal form and definition of } \underline{step}) \\ & \quad a + \sum_{i=1}^m \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\ & \quad + \sum_{i=1}^n \underline{rel}(z_i, nf(c\ (sub\ e'_1\ z_1 \dots z_n) \dots \\ & \quad \quad \quad (sub\ e'_m\ z_1 \dots z_n))) \cdot \underline{step}(e_i) \\ &= \text{(by Observation 5 because } \underline{tree}(e) = \underline{tree}(e[z_i \rightsquigarrow z_i]) \text{)} \end{aligned}$$

$$\begin{aligned}
& a + \sum_{i=1}^m \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf((c\ (sub\ e'_1\ z_1 \dots z_n) \dots \\
& \quad \quad \quad (sub\ e'_m\ z_1 \dots z_n))[z_i \rightsquigarrow z_i])) \cdot \underline{step}(e_j) \\
= & \quad (\text{by reduction}) \\
& a + \sum_{i=1}^m \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\
& \quad + \sum_{i=1}^n \underline{rel}(z_i, nf(sub\ (c\ e'_1 \dots e'_m)\ z_1 \dots z_n)) \cdot \underline{step}(e_i) \\
= & \quad (\text{by induction hypothesis [} m \text{ times]}) \\
& a + \sum_{j=1}^m \left(\underline{step}(e'_j) + \sum_{l=1}^n \underline{rel}(z_l, e'_j[\Pi_i/z_i]) \cdot \underline{step}(z_l) \right) + \\
& \quad + \sum_{i=1}^n \underline{rel}(z_i, nf(sub\ (c\ e'_1 \dots e'_m)\ z_1 \dots z_n)) \cdot \underline{step}(e_i) \\
= & \quad (\text{by definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0) \\
& a + \sum_{i=1}^m \underline{step}(e'_i) + \\
& \quad + \sum_{i=1}^n \underline{rel}(z_i, nf(sub\ (c\ e'_1 \dots e'_m)\ z_1 \dots z_n)) \cdot \underline{step}(e_i) \\
= & \quad (\text{by Lemma 1, } \underline{tree}(e)[\Pi_i/z_i] = \underline{tree}(e[\Pi_i/z_i]), \text{ and Observation 5}) \\
& a + \sum_{i=1}^m \underline{step}(e'_i) + \\
& \quad + \sum_{i=1}^n \underline{rel}(z_i, (c\ e'_1 \dots e'_m)[\Pi_i/z_i]) \cdot \underline{step}(e_i) \\
= & \quad (\text{by definition of } \underline{step}) \\
& \underline{step}(c\ e'_1 \dots e'_m) + \sum_{i=1}^n \underline{rel}(z_i, (c\ e'_1 \dots e'_m)[\Pi_i/z_i]) \cdot \underline{step}(e_i)
\end{aligned}$$

(iii) Finally, let $e_0 = e'_0[z_i \rightsquigarrow e'_i]$ for some $e'_0 \in E_{C^0, n}(Z_n)$ and $e'_0, \dots, e'_n \in E_{C^0, n}$.

$$\begin{aligned}
& \underline{step}(nf(sub\ e'_0[z_i \rightsquigarrow e'_i]\ e_1 \dots e_n)) \\
= & \quad (\text{by reduction}) \\
& \underline{step}(nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n][z_i \rightsquigarrow e_i])) \\
= & \quad (\text{by normal form because } e_i \text{ is in normal form}) \\
& \underline{step}(nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n])[z_i \rightsquigarrow e_i]) \\
= & \quad (\text{by definition of } \underline{step}) \\
& \underline{step}(nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n])) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n])) \cdot \underline{step}(e_j) \\
= & \quad (\text{by normal form and definition of } \underline{step} \text{ because } e'_0 \text{ is in normal form}) \\
& \underline{step}(e'_0) + \sum_{i=1}^n \underline{rel}(z_i, e'_0) \cdot \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n])) \cdot \underline{step}(e_j) \\
= & \quad (\text{by Observation 5 because } \underline{tree}(e) = \underline{tree}(e[z_i \rightsquigarrow z_i])) \\
& \underline{step}(e'_0) + \sum_{i=1}^n \underline{rel}(z_i, e'_0) \cdot \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(e'_0[z_i \rightsquigarrow sub\ e'_i\ z_1 \dots z_n][z_i \rightsquigarrow z_i])) \cdot \underline{step}(e_j) \\
= & \quad (\text{by reduction})
\end{aligned}$$

$$\begin{aligned}
& \underline{step}(e'_0) + \sum_{i=1}^n \underline{rel}(z_i, e'_0) \cdot \underline{step}(nf(sub\ e'_i\ z_1 \dots z_n)) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(sub\ e'_0[z_i \rightsquigarrow e'_i]\ z_1 \dots z_n)) \cdot \underline{step}(e_j) \\
= & \quad (\text{by induction hypothesis [n times]}) \\
& \underline{step}(e'_0) + \sum_{j=1}^n \underline{rel}(z_j, e'_0) \cdot \\
& \quad \cdot \left(\underline{step}(e'_j) + \sum_{l=1}^n \underline{rel}(z_l, e'_j[\Pi_l/z_l]) \cdot \underline{step}(z_l) \right) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(sub\ e'_0[z_i \rightsquigarrow e'_i]\ z_1 \dots z_n)) \cdot \underline{step}(e_j) \\
= & \quad (\text{by definition of } \underline{step} \text{ because } \underline{step}(z_l) = 0) \\
& \underline{step}(e'_0) + \sum_{i=1}^n \underline{rel}(z_i, e'_0) \cdot \underline{step}(e'_i) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf(sub\ e'_0[z_i \rightsquigarrow e'_i]\ z_1 \dots z_n)) \cdot \underline{step}(e_j) \\
= & \quad (\text{by Lemma 1, } \underline{tree}(e)[\Pi_i/z_i] = \underline{tree}(e[\Pi_i/z_i]), \text{ and Observation 5}) \\
& \underline{step}(e'_0) + \sum_{i=1}^n \underline{rel}(z_i, e'_0) \cdot \underline{step}(e'_i) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, e'_0[z_i \rightsquigarrow e'_i][\Pi_i/z_i]) \cdot \underline{step}(e_j) \\
= & \quad (\text{by definition of } \underline{step}) \\
& \underline{step}(e'_0[z_i \rightsquigarrow e'_i]) + \sum_{j=1}^n \underline{rel}(z_j, e'_0[z_i \rightsquigarrow e'_i][\Pi_i/z_i]) \cdot \underline{step}(e_j)
\end{aligned}$$

This ends the proof of the statement. \square

Lemma 7. Let $M = (F, Sub, C, \Pi, R_1, R_2, r_{in})$ with $Sub = \{sub^{(n+1)}\}$ be an n -sntt, and let $acc(M) = (acc(F), C, acc(R_1), acc(r_{in}))$ be the n -satt constructed from M by accumulation.

(a) Let $t \in T_C$ and $e_1, \dots, e_n \in E_{C^\circ, n}$.

$$\begin{aligned}
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, f\ t\ e_1 \dots e_n)) \\
= & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, (f\ t\ z_1 \dots z_n)[z_i \rightsquigarrow e_i]))
\end{aligned}$$

(b) Let $r \in RHS(F, Sub, C, X_k)$, $t_1, \dots, t_k \in T_C$, and $e_1, \dots, e_n \in E_{C^\circ, n}$ for some $k \in \mathbb{N}$.

$$\begin{aligned}
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub}\ r\ e_1 \dots e_n)[x_i/t_i])) \\
= & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub}\ r\ z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i]))
\end{aligned}$$

Proof. We prove the two statements by simultaneous induction. We abbreviate $nf(\Rightarrow_{acc(R_1)^\circ}, e)$ simply by $nf_1(e)$ and similarly $nf(\Rightarrow_{\overline{R_2}}, e)$ by $nf_2(e)$. By induction on t we first show Statement (a) with the help of Statement (b). Let $t = c\ t_1 \dots t_k$ for some $k \in \mathbb{N}$, $c \in C^{(k)}$, and $t_1, \dots, t_k \in T_C$.

$$\begin{aligned}
& \underline{step}(nf_1(f(c\ t_1 \dots t_k)\ e_1 \dots e_n)) \\
= & \quad (\text{by reduction and definition of } acc(R_1)^\circ) \\
& \underline{step}(nf_1((\circ\ rhs_{acc(R_1), f, c})[x_i/t_i][y_i/e_i])) \\
= & \quad (\text{by } nf_1, \text{ definition of } \underline{step} \text{ and } acc(R_1), \text{ and substitution})
\end{aligned}$$

$$\begin{aligned}
& 1 + \underline{step}(nf_1(nf_2(\underline{sub} \text{ rhs}_{R_1, f, c} y_1 \dots y_n)[x_i/t_i] [y_i/e_i])) \\
= & \text{ (by substitution because } e_i \text{ is in normal form)} \\
& 1 + \underline{step}(nf_1(nf_2(\underline{sub} \text{ rhs}_{R_1, f, c} e_1 \dots e_n)[x_i/t_i])) \\
= & \text{ (by Statement (b))} \\
& 1 + \underline{step}(nf_1(nf_2(\underline{sub} \text{ rhs}_{R_1, f, c} z_1 \dots z_n)[x_i/t_i] [z_i \rightsquigarrow e_i])) \\
= & \text{ (by } nf_1, \text{ definition of } \underline{step}, \text{ and substitution)} \\
& \underline{step}(nf_1((\circ nf_2(\underline{sub} \text{ rhs}_{R_1, f, c} y_1 \dots y_n))[x_i/t_i] [y_i/z_i] [z_i \rightsquigarrow e_i])) \\
= & \text{ (by definition of } acc(R_1)) \\
& \underline{step}(nf_1((\circ \text{ rhs}_{acc(R_1), f, c}[x_i/t_i] [y_i/z_i] [z_i \rightsquigarrow e_i])) \\
= & \text{ (by reduction and definition of } acc(R_1)^\circ) \\
& \underline{step}(nf_1((f (c t_1 \dots t_k) z_1 \dots z_n)[z_i \rightsquigarrow e_i]))
\end{aligned}$$

Now we prove Statement (b) by induction on r .

(i) First, let $r = II_j$ for some $j \in [n]$.

$$\begin{aligned}
& \underline{step}(nf_1(nf_2(\underline{sub} II_j e_1 \dots e_n)[x_i/t_i])) \\
= & \text{ (by reduction, substitution, and } nf_1 \text{ because } e_j \text{ is in normal form)} \\
& \underline{step}(e_j) \\
= & \text{ (by definition of } \underline{step} \text{ and } \underline{rel}) \\
& \underline{step}(z_j) + \sum_{i=1}^n \underline{rel}(z_i, z_j) \cdot \underline{step}(e_i) \\
= & \text{ (by definition of } \underline{step}) \\
& \underline{step}(z_j[z_i \rightsquigarrow e_i]) \\
= & \text{ (by } nf_1, \text{ substitution, and reduction)} \\
& \underline{step}(nf_1(nf_2(\underline{sub} II_j z_1 \dots z_n)[x_i/t_i] [z_i \rightsquigarrow e_i]))
\end{aligned}$$

(ii) Second, let $r = f x_j$ for some $f \in F$ and $j \in [k]$.

$$\begin{aligned}
& \underline{step}(nf_1(nf_2(\underline{sub} (f x_j) e_1 \dots e_n)[x_i/t_i])) \\
= & \text{ (by reduction)} \\
& \underline{step}(nf_1((f x_j e_1 \dots e_n)[x_i/t_i])) \\
= & \text{ (by substitution)} \\
& \underline{step}(nf_1(f t_j e_1 \dots e_n)) \\
= & \text{ (by Statement (a))} \\
& \underline{step}(nf_1((f t_j z_1 \dots z_n)[z_i \rightsquigarrow e_i])) \\
= & \text{ (by substitution)} \\
& \underline{step}(nf_1((f x_j z_1 \dots z_n)[x_i/t_i] [z_i \rightsquigarrow e_i])) \\
= & \text{ (by reduction)}
\end{aligned}$$

$$\underline{step}(nf_1(nf_2(\underline{sub}(f x_j) z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i]))$$

(iii) Third, let $r = c r_1 \dots r_m$ for some $m \in \mathbb{N}$, $c \in (C - \Pi)^{(m)}$, and $r_1, \dots, r_m \in RHS(F, Sub, C, X_k)$.

$$\begin{aligned} & \underline{step}(nf_1(nf_2(\underline{sub}(c r_1 \dots r_m) e_1 \dots e_n)[x_i/t_i])) \\ = & \text{ (by reduction)} \\ & \underline{step}(nf_1(nf_2((c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[z_i \rightsquigarrow e_i])[x_i/t_i])) \\ = & \text{ (by } nf_1, \text{ substitution, and } nf_2) \\ & \underline{step}(nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i][z_i \rightsquigarrow e_i])) \\ = & \text{ (by definition of } \underline{step}) \\ & \underline{step}(nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) + \\ & + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(e_j) \\ = & \text{ (by definition of } \underline{step} \text{ because } \underline{step}(z_j) = 0) \\ & \underline{step}(nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) + \\ & + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(e_j) + \\ & + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(z_j) \\ = & \text{ (by definition of } \underline{step} \text{ and Observation 5)} \\ & \underline{step}(nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i][z_i \rightsquigarrow z_i])) + \\ & + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i][z_i \rightsquigarrow z_i])) \cdot \underline{step}(e_j) \\ = & \text{ (by definition of } \underline{step}) \\ & \underline{step}(nf_1(nf_2(c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i][z_i \rightsquigarrow z_i][z_i \rightsquigarrow e_i])) \\ = & \text{ (by } nf_1, \text{ substitution, and } nf_2) \\ & \underline{step}(nf_1(nf_2((c (\underline{sub} r_1 z_1 \dots z_n) \dots \\ & \quad (\underline{sub} r_m z_1 \dots z_n))[z_i \rightsquigarrow z_i])[x_i/t_i][z_i \rightsquigarrow e_i])) \\ = & \text{ (by reduction)} \\ & \underline{step}(nf_1(nf_2(\underline{sub}(c r_1 \dots r_m) z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i])) \\ = & \text{ (by } nf_1) \end{aligned}$$

$$\begin{aligned}
& \underline{step}(nf_1(nf_2(\underline{sub} (c r_1 \dots r_m) z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i])) \\
\text{(iv) Finally, let } r &= \underline{sub} r_0 \dots r_n \text{ for some } r_0, \dots, r_n \in RHS(F, Sub, C, X_k). \\
& \underline{step}(nf_1(nf_2(\underline{sub} (sub r_0 \dots r_n) e_1 \dots e_n)[x_i/t_i])) \\
&= \text{ (by reduction)} \\
& \underline{step}(nf_1(nf_2((\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_n z_1 \dots z_n))[z_i \rightsquigarrow e_i])[x_i/t_i])) \\
&= \text{ (by } nf_1, \text{ substitution, and } nf_2) \\
& \underline{step}(nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])[z_i \rightsquigarrow e_i]) \\
&= \text{ (by definition of } \underline{step}) \\
& \underline{step}(nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(e_j) \\
&= \text{ (by definition of } \underline{step} \text{ because } \underline{step}(z_j) = 0) \\
& \underline{step}(nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(e_j) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])) \cdot \underline{step}(z_j) \\
&= \text{ (by definition of } \underline{step} \text{ and Observation 5)} \\
& \underline{step}(nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])[z_i \rightsquigarrow z_i]) + \\
& \quad + \sum_{j=1}^n \underline{rel}(z_j, nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])[z_i \rightsquigarrow z_i])) \cdot \underline{step}(e_j) \\
&= \text{ (by definition of } \underline{step}) \\
& \underline{step}(nf_1(nf_2(\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_m z_1 \dots z_n))[x_i/t_i])[z_i \rightsquigarrow z_i][z_i \rightsquigarrow e_i]) \\
&= \text{ (by } nf_1, \text{ substitution, and } nf_2) \\
& \underline{step}(nf_1(nf_2((\underline{sub} r_0 (\underline{sub} r_1 z_1 \dots z_n) \dots \\
& \quad (\underline{sub} r_m z_1 \dots z_n))[z_i \rightsquigarrow z_i])[x_i/t_i])[z_i \rightsquigarrow e_i]) \\
&= \text{ (by reduction)} \\
& \underline{step}(nf_1(nf_2(\underline{sub} (sub r_0 \dots r_n) z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i])) \\
&= \text{ (by } nf_1) \\
& \underline{step}(nf_1(nf_2(\underline{sub} (sub r_0 \dots r_n) z_1 \dots z_n)[x_i/t_i][z_i \rightsquigarrow e_i]))
\end{aligned}$$

Thus the proof is complete. \square

Lemma 8 (see Lemma 20). Let $M = (F, Sub, C, \Pi, R_1, R_2, r_{in})$ with $Sub = \{sub^{(n+1)}\}$ be an n -sntt, and let $acc(M) = (acc(F), C, acc(R_1), acc(r_{in}))$ be the n -satt constructed from M by accumulation.

(a) For every $f \in F$ and $t \in T_C$:

$$\underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub(f t) z_1 \dots z_n)) = \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, f t z_1 \dots z_n)) .$$

(b) For every $k \in \mathbb{N}$, $r \in RHS(F, Sub, C, X_k)$, and $t_1, \dots, t_k \in T_C$:

$$\begin{aligned} & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub r[x_i/t_i] z_1 \dots z_n)) \\ &= \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r z_1 \dots z_n)[x_i/t_i])) . \end{aligned}$$

Proof. We prove the two statements by simultaneous induction. Statement (a) is proved by structural induction on t , so let $t = c t_1 \dots t_k$ for some $k \in \mathbb{N}$, $c \in C^{(k)}$, and $t_1, \dots, t_k \in T_C$.

$$\begin{aligned} & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub(f(c t_1 \dots t_k)) z_1 \dots z_n)) \\ &= \text{(by reduction and definition of } rhs_{R_1, f, c}) \\ & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub(\circ rhs_{R_1, f, c}[x_i/t_i] z_1 \dots z_n))) \\ &= \text{(by substitution and reduction)} \\ & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, (\circ(sub rhs_{R_1, f, c}[x_i/t_i] z_1 \dots z_n))[z_i \rightsquigarrow z_i])) \\ &= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot) \text{ and definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0) \\ & 1 + \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub rhs_{R_1, f, c}[x_i/t_i] z_1 \dots z_n)) \\ &= \text{(by Statement (b))} \\ & 1 + \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} rhs_{R_1, f, c} z_1 \dots z_n)[x_i/t_i])) \\ &= \text{(by substitution, } nf(\Rightarrow_{acc(R_1)^\circ}, \cdot), \text{ and definition of } \underline{step}) \\ & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, (\circ(nf(\Rightarrow_{\overline{R_2}}, \underline{sub} rhs_{R_1, f, c} y_1 \dots y_n)))[x_i/t_i] [y_i/z_i])) \\ &= \text{(by reduction and definition of } rhs_{acc(R_1)^\circ, f, c} \text{ and } rhs_{acc(R_1)^\circ, f, c}) \\ & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, f(c t_1 \dots t_k) z_1 \dots z_n)) \end{aligned}$$

Statement (b) is proved by induction on r .

(i) First let $r = \Pi_j$ for some $j \in [n]$.

$$\begin{aligned} & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, sub \Pi_j[x_i/t_i] z_1 \dots z_n)) \\ &= \text{(by substitution and reduction)} \\ & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, z_j)) \\ &= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot), nf(\Rightarrow_{acc(R_1)^\circ}, \cdot), \text{ and substitution)} \\ & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, z_j[x_i/t_i])) \end{aligned}$$

$$\begin{aligned}
&= \text{(by reduction)} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} \Pi_j z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

(ii) Now let $r = f x_j$ for some $f \in F$ and $j \in [k]$.

$$\begin{aligned}
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} (f x_j)[x_i/t_i] z_1 \dots z_n)) \\
&= \text{(by substitution)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} (f t_j) z_1 \dots z_n)) \\
&= \text{(by Statement (a))} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, f t_j z_1 \dots z_n)) \\
&= \text{(by substitution)} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, (f x_j z_1 \dots z_n)[x_i/t_i])) \\
&= \text{(by reduction)} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} (f x_j) z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

(iii) Let $r = c r_1 \dots r_m$ for some $m \in \mathbb{N}$, $c \in (C - \Pi)^{(m)}$, and $r_1, \dots, r_m \in RHS(F, Sub, C, X_k)$. Moreover, let $a = 1$ if $c = \circ$ and $a = 0$ otherwise.

$$\begin{aligned}
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} (c r_1 \dots r_m)[x_i/t_i] z_1 \dots z_n)) \\
&= \text{(by substitution)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} (c r_1[x_i/t_i] \dots r_m[x_i/t_i]) z_1 \dots z_n)) \\
&= \text{(by reduction)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, (c (\underline{sub} r_1[x_i/t_i] z_1 \dots z_n) \dots \\
&\quad \quad \quad (\underline{sub} r_m[x_i/t_i] z_1 \dots z_n))[z_i \rightsquigarrow z_i])) \\
&= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot) \text{ and definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0) \\
&\quad a + \sum_{j=1}^m \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_j[x_i/t_i] z_1 \dots z_n)) \\
&= \text{(by induction hypothesis [} m \text{ times])} \\
&\quad a + \sum_{j=1}^m \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])) \\
&= \text{(by } nf(\Rightarrow_{acc(R_1)^\circ}, \cdot), nf(\Rightarrow_{\overline{R_2}}, \cdot), \text{ substitution, and definition of } \underline{step}) \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \\
&\quad \quad \quad c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n)[x_i/t_i])) \\
&= \text{(by definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0) \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \\
&\quad \quad \quad c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n)[x_i/t_i])[z_i \rightsquigarrow z_i]) \\
&= \text{(by } nf(\Rightarrow_{acc(R_1)^\circ}, \cdot), nf(\Rightarrow_{\overline{R_2}}, \cdot), \text{ and substitution)} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \\
&\quad \quad \quad (c (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_m z_1 \dots z_n))[z_i \rightsquigarrow z_i])[x_i/t_i]))
\end{aligned}$$

$$\begin{aligned}
&= \text{(by reduction)} \\
&\quad \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub}(c r_1 \dots r_m) z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

(iv) Finally, let $r = \underline{sub} r_0 \dots r_n$ for some $r_0, \dots, r_n \in RHS(F, Sub, C, X_k)$.

$$\begin{aligned}
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub}(\underline{sub} r_0 \dots r_n)[x_i/t_i] z_1 \dots z_n)) \\
&= \text{(by substitution)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub}(r_0[x_i/t_i] \dots r_n[x_i/t_i]) z_1 \dots z_n)) \\
&= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot) \text{)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \\
&\quad \quad \underline{sub} nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i]) z_1 \dots z_n)) \\
&= \text{(by Lemma 6)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i])) + \\
&\quad + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i])[II_i/z_i]) \cdot \underline{step}(z_j) \\
&= \text{(because } \underline{step}(z_j) = 0 \text{)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] \dots r_n[x_i/t_i])) \\
&= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot) \text{)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i]) \dots \\
&\quad \quad \quad nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_n[x_i/t_i]))) \\
&= \text{(by Lemma 6)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
&\quad + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \\
&\quad \quad \cdot \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_j[x_i/t_i])) \\
&= \text{(by definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0 \text{)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
&\quad + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \\
&\quad \quad \cdot \left(\underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_j[x_i/t_i])) + \right. \\
&\quad \quad \quad \left. + \sum_{l=1}^n \underline{rel}(z_l, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_j[x_i/t_i])[II_i/z_i]) \cdot \underline{step}(z_l) \right) \\
&= \text{(by Lemma 6)} \\
&\quad \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
&\quad + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \\
&\quad \quad \cdot \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_j[x_i/t_i]) z_1 \dots z_n)) \\
&= \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot) \text{)}
\end{aligned}$$

$$\begin{aligned}
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_j[x_i/t_i] z_1 \dots z_n)) \\
= & \text{(by Lemma 1 and Observation 5 and } \underline{step}(z_j) = 0) \\
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \underline{step}(z_j) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] z_1 \dots z_n)) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_j[x_i/t_i] z_1 \dots z_n)) \\
= & \text{(by induction hypothesis [n times])} \\
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i])[II_i/z_i]) \cdot \underline{step}(z_j) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] z_1 \dots z_n)) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by Lemmata 6 and 3 and Observation 5)} \\
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_0[x_i/t_i] z_1 \dots z_n))) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by } nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \cdot)) \\
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_0[x_i/t_i] z_1 \dots z_n)) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by induction hypothesis)} \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) + \\
& + \sum_{j=1}^n \underline{rel}(z_j, nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i])) \cdot \\
& \quad \cdot \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])) \\
= & \text{(by definition of } \underline{step}) \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i]) \\
& \quad [z_j \rightsquigarrow nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])]) \\
= & \text{(by } nf(\Rightarrow_{acc(R_1)^\circ}, \cdot)) \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 z_1 \dots z_n)[x_i/t_i] \\
& \quad [z_j \rightsquigarrow nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_j z_1 \dots z_n)[x_i/t_i])]) \\
= & \text{(by Lemma 7)}
\end{aligned}$$

$$\begin{aligned}
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_1 z_1 \dots z_n)[x_i/t_i] \dots \\
& \quad nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_n z_1 \dots z_n)[x_i/t_i])[x_i/t_i])) \\
= & \text{ (by } nf(\Rightarrow_{\overline{R_2}}, \cdot) \text{ and substitution)} \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[x_i/t_i])) \\
= & \text{ (by definition of } \underline{step} \text{ because } \underline{step}(z_i) = 0) \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[x_i/t_i])[z_i \rightsquigarrow z_i]) \\
= & \text{ (by } nf(\Rightarrow_{acc(R_1)^\circ}, \cdot), \text{ substitution, and } nf(\Rightarrow_{\overline{R_2}}, \cdot)) \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, (\underline{sub} r_0 \\
& \quad (\underline{sub} r_1 z_1 \dots z_n) \dots (\underline{sub} r_n z_1 \dots z_n))[z_i \rightsquigarrow z_i])[x_i/t_i])) \\
= & \text{ (by reduction)} \\
& \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} (sub r_0 \dots r_n) z_1 \dots z_n)[x_i/t_i]))
\end{aligned}$$

This concludes the proof. \square

Theorem 9 (see Theorem 21). Let $M = (F, \{sub^{(n+1)}\}, C, \Pi, R_1, R_2, r_{in})$ be an n -sntt and $acc(M) = (acc(F), C, acc(R_1), acc(r_{in}))$. Then, for every $t \in T_C$:

$$\underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_{in}[x_1/t])) = \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, acc(r_{in})[x_1/t])) .$$

Proof.

$$\begin{aligned}
& \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_{in}[x_1/t])) \\
= & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_{in}[x_1/t] z_1 \dots z_n))) \quad (*) \\
= & \underline{step}(nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, \underline{sub} r_{in}[x_1/t] z_1 \dots z_n)) \quad (\text{confluence}) \\
= & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_{in} z_1 \dots z_n)[x_1/t])) \quad (\text{Lm. 8(b)}) \\
= & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, nf(\Rightarrow_{\overline{R_2}}, \underline{sub} r_{in} \Pi_1 \dots \Pi_n)[x_1/t])) \quad (**) \\
= & \underline{step}(nf(\Rightarrow_{acc(R_1)^\circ}, acc(r_{in})[x_1/t])) \quad (\text{Def. } acc(r_{in}))
\end{aligned}$$

The statement (*) holds, because sub neither deletes nor copies \circ -symbols and preserves sharings in its recursion argument (in particular because of the last rule of Definition 6). (*) can be formally proved by induction on $nf(\Rightarrow_{R_1^\circ \cup R_2^\circ}, r_{in}[x_1/t])$. Note that (*) does not hold without the outer application of \underline{step} .

The statement (**) holds, because for every $j \in [n]$, both, z_j and Π_j do not contain \circ -symbols and are normal forms w.r.t. $\Rightarrow_{\overline{R_2}}$ and $\Rightarrow_{acc(R_1)^\circ}$. (**) can be formally proved by induction on r_{in} . \square