

# Parsing Algorithms Based on Tree Automata

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# Parsing and CFG

## Example (Context-free grammar)

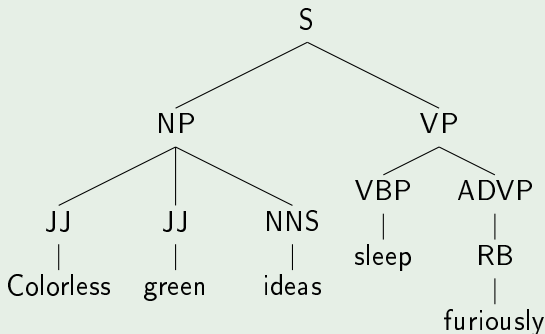
|                                     |                                  |
|-------------------------------------|----------------------------------|
| $S \rightarrow NP VP$               | $NP \rightarrow JJ JJ NNS$       |
| $VP \rightarrow VBP ADVP$           | $ADVP \rightarrow RB$            |
| $JJ \rightarrow \textit{Colorless}$ | $JJ \rightarrow \textit{green}$  |
| $NNS \rightarrow \textit{ideas}$    | $VBP \rightarrow \textit{sleep}$ |
| $RB \rightarrow \textit{furiously}$ |                                  |

## Derivation

$S \rightarrow^* \textit{Colorless green ideas sleep furiously}$

# Parse tree

## Example

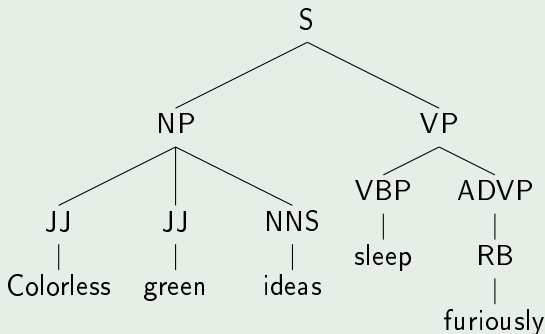


## Remark

We are interested in the parse tree, not just whether  $S \rightarrow^* w!$

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# Local tree language

## Definition

A **local tree grammar** is a grammar with rules of the form

$$S \rightarrow \begin{array}{c} S \\ \swarrow \quad | \quad \searrow \\ N_1 \quad \cdots \quad N_k \end{array}$$

The such generated languages are the **local tree languages**.

## Theorem

*The derivations of a context-free grammar are a local tree language.*

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Regular tree grammars are local tree grammars with hidden states.

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# Back to parsing

## Observation

Most CFG-parsers are regular tree grammars (+ control) because

- they are based on a CFG ( $\rightarrow$  local tree grammar) and
- have hidden states (or features)

## Alternative

The features can be made explicit in the parse tree structure.

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## Observation

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# Regular restriction

## Theorem [Bar-Hillel et al '64]

The intersection of a context-free language with a regular language is again context-free.

## Regular restriction — Trivial approach

### Theorem

*For every regular language  $L$ , the set of all trees, whose yield is in  $L$ , is regular.*

### Theorem

*The intersection of two regular tree languages is regular.*

### Theorem

*For every regular tree language the restriction to a regular language of yields is again regular.*

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# The End?

Thank you for your attention!

# Contents

- 1 Motivation
- 2 Weighted Tree Grammars
- 3 Bar-Hillel Construction
- 4 Further Topics



# Weighted tree grammar

## Definition

A weighted tree grammar is a structure  $(\mathcal{N}, \Sigma, \mathbb{K}, N_0, R)$  with rules of the form

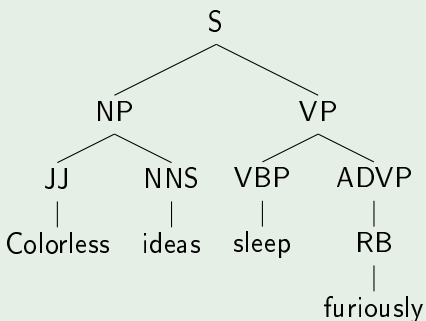
$$N \xrightarrow{c} \begin{array}{c} V \\ \swarrow \quad | \quad \searrow \\ N_1 \quad \dots \quad N_k \end{array}$$

where

- $N, N_1, \dots, N_k \in \mathcal{N}$  are nonterminals
- $c \in \mathbb{K}$  is a weight (taken from a semiring)
- $V \in \Sigma$  is a terminal symbol

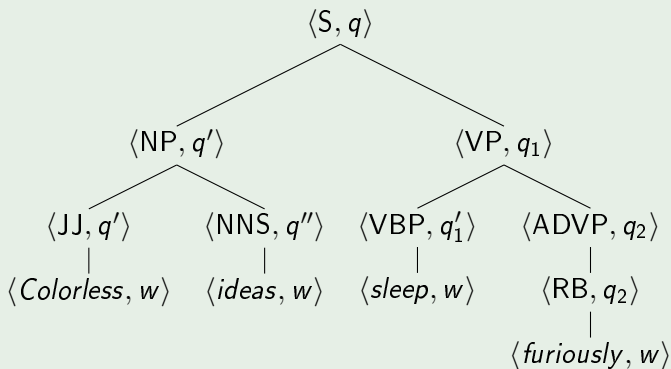
## Runs

## Example (Input tree)



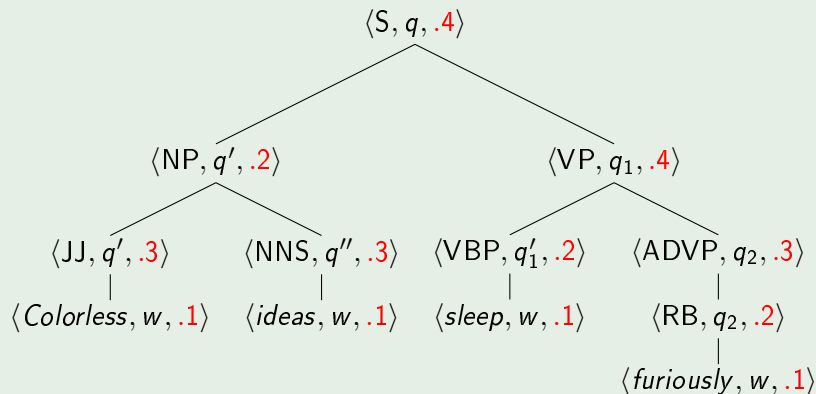
## Runs

## Example (Run)



## Runs

## Example (Run with weights)



# Weight of a run

## Definition

The weight of a run is obtained by multiplying the weights in it.

## Definition

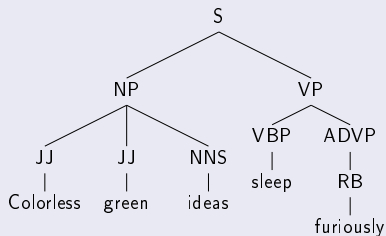
The weight of an input tree is obtained by adding the weights of all runs on it.

# Contents

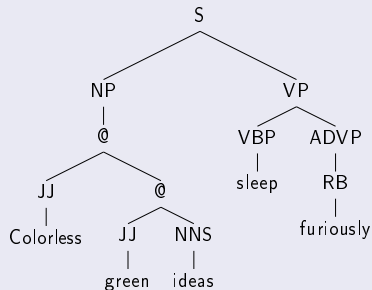
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# Binarization

## Input tree



## Binarized tree

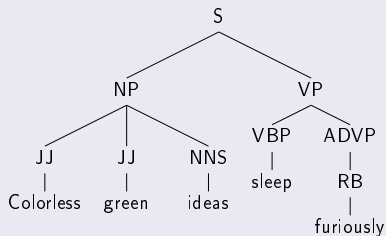


## Theorem

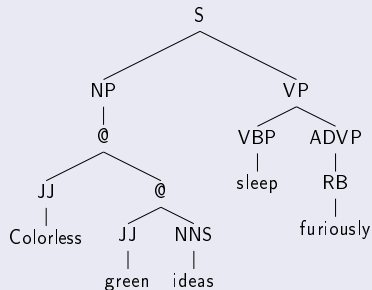
A tree language is regular if and only if its binarization is regular (also holds in the weighted setting).

# Binarization

## Input tree



## Binarized tree



## Theorem

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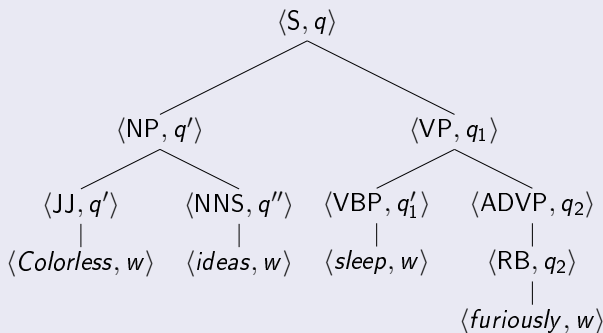


# Individual runs

## Run on the yield

$(p)$  Colorless  $(p_1)$  ideas  $(p_2)$  sleep  $(p_3)$  furiously  $(p')$

## Run on the input tree

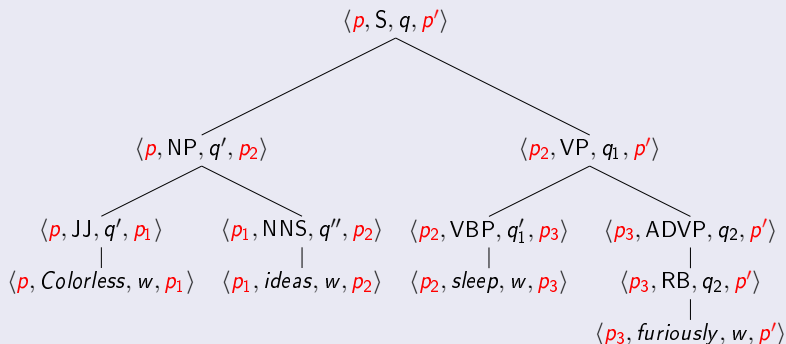


# Bar-Hillel construction

## Run on the yield

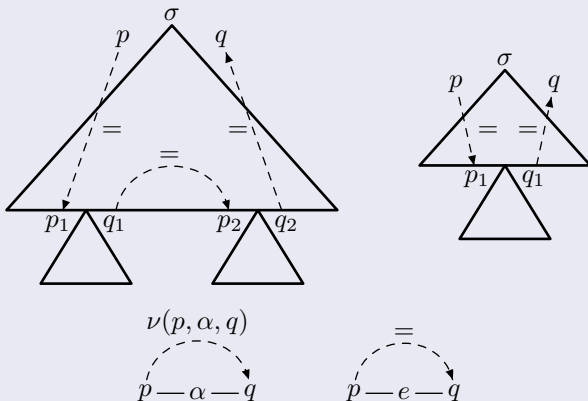
$\langle p \rangle$  Colorless  $\langle p_1 \rangle$  ideas  $\langle p_2 \rangle$  sleep  $\langle p_3 \rangle$  furiously  $\langle p' \rangle$

## Composite run



# Bar-Hillel construction (cont'd)

## Illustration



## Bar-Hillel construction (cont'd)

### Theorem

*The regular restriction of a regular tree language is regular (also in the weighted setting).*

### Remark

Complexity:  $O(mn^3)$

- $m$ : size of the regular tree grammar
- $n$ : size of the regular grammar (or input string)

### Conclusion

We can parse with regular tree grammars in  $O(mn^3)$ .

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## Further topics — Probability mass

### Definition

The probability mass of a nonterminal is the sum of the weights of runs with that nonterminal at the root.

### Algorithm

The probability mass of a state can be computed using fix-point iteration.

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## Further topics — Normalization

### Definition

A weighted regular tree grammar is

- **convergent** if the sum of weights assigned to trees are (uniformly) bounded
- **proper** if the probabilities of rules for one nonterminal add to 1
- **consistent** if it computes a probability distribution.

### Theorem

For every convergent grammar, there exist a scaled proper and consistent grammar.



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For every convergent grammar, there exist a scaled proper and consistent grammar.

## Further topics — Best parse

### Theorem

*For an unambiguous grammar we can compute the best parse in linear time.*

### Proof.

Using a variant of Knuth's algorithm.

## References

- [Bar-Hillel, Perles, Shamir](#): *On formal properties of simple phrase structure grammars*. Language and Information, 1964
- [Berstel, Reutenauer](#): Recognizable formal power series on trees. *Theoret. Comput. Sci.* 18, p. 115–148, 1982
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- [Nederhof, Satta](#): Probabilistic parsing as intersection. In IWPT 2003

Thank you for your attention!